

Some recent improvements of parallel-in-time algorithms

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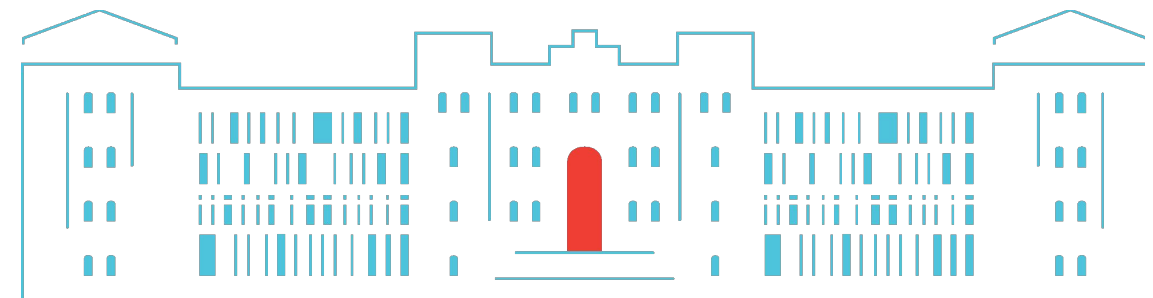
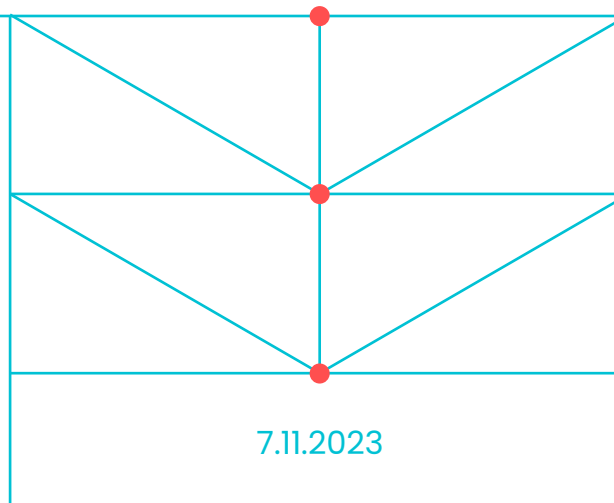
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of Education
and Research**



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Daniel Ruprecht

Results by Abdul Ibrahim, Thibaut Lunet, Thomas Baumann

The problem with parallelizing only in space

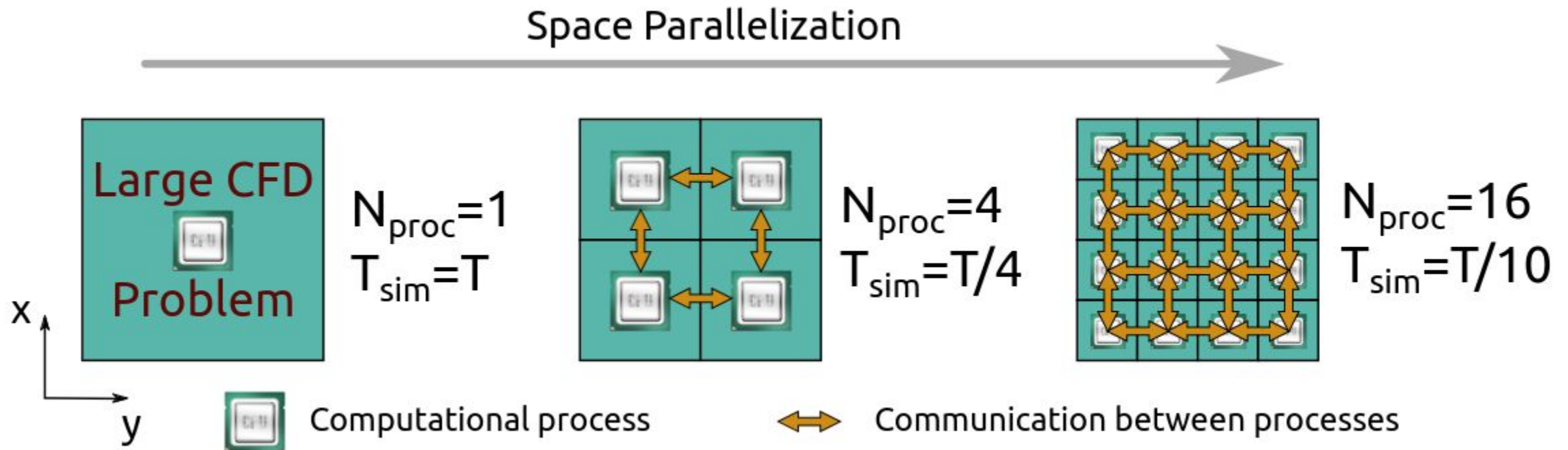
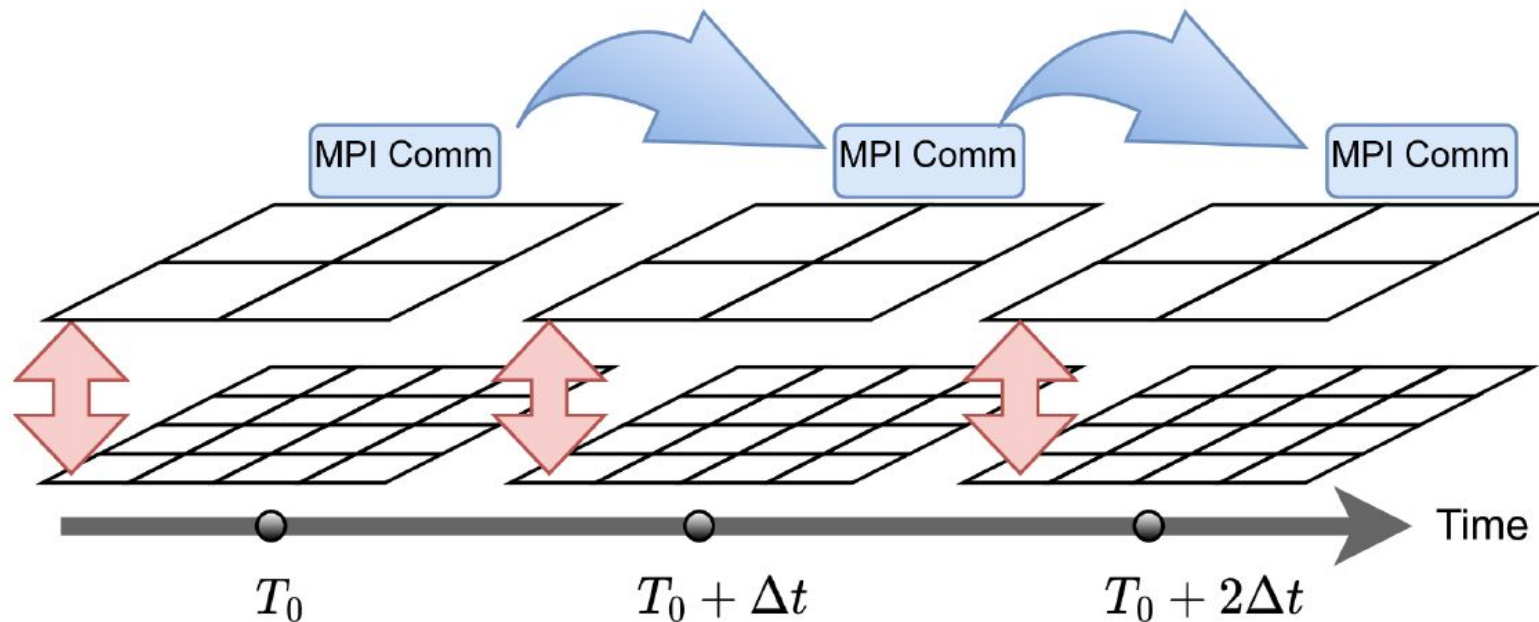


Figure courtesy T. Lunet

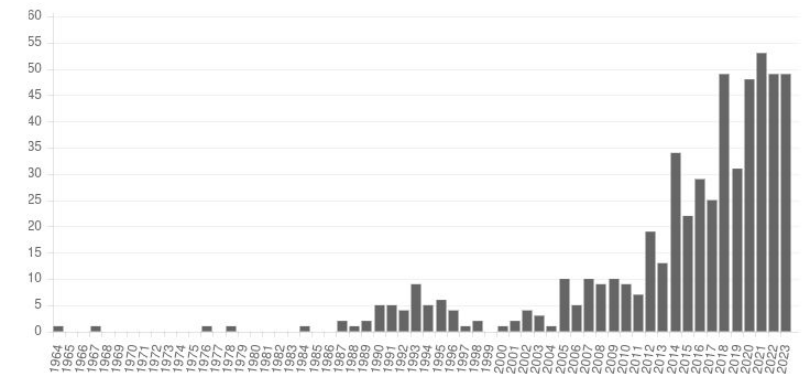
- Spatial strong scaling eventually saturates
- Even with perfect weak scaling in space, increased time resolution still increases solution times

Parallel-across-the-steps: the basic idea



We cannot completely avoid serial dependency in time ... but we can *relax* it.

- First idea published in 1964
- Small surge of interest in 1990's
- rapid growth of the field since 2001



What PinT promises to deliver

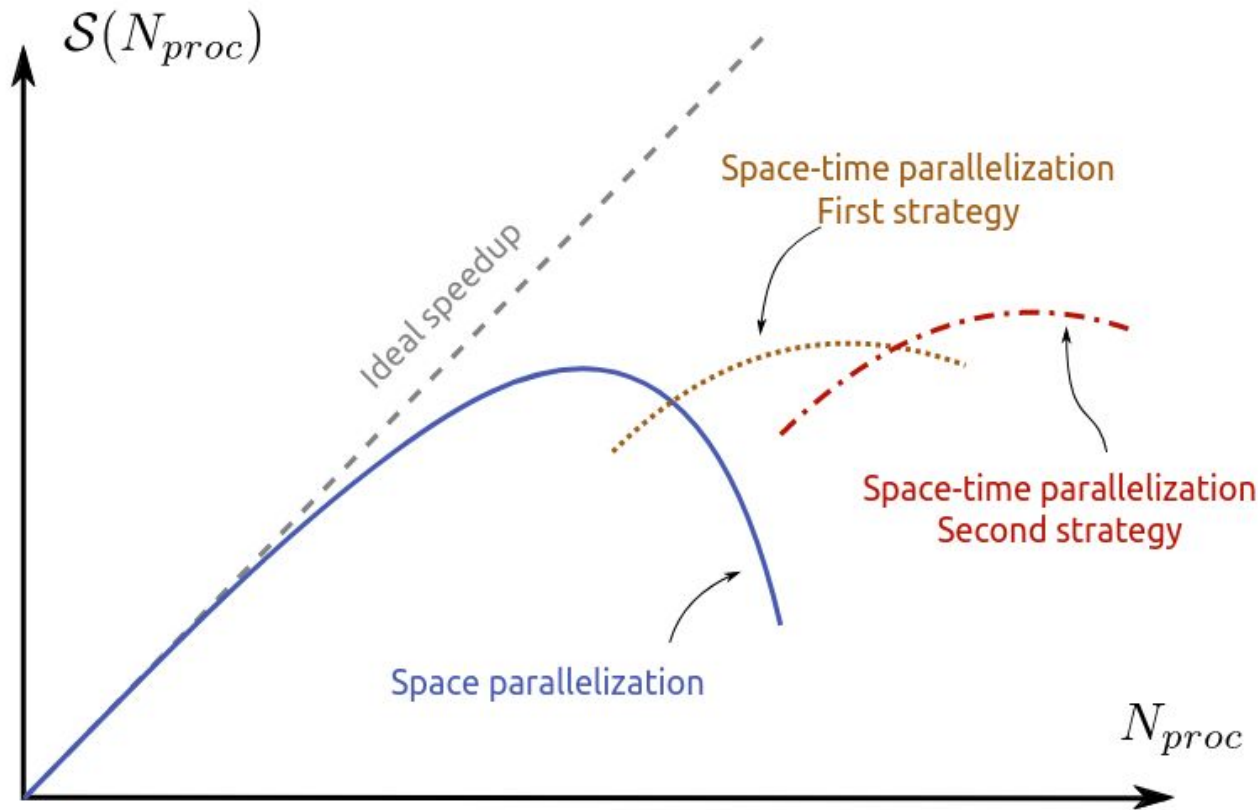
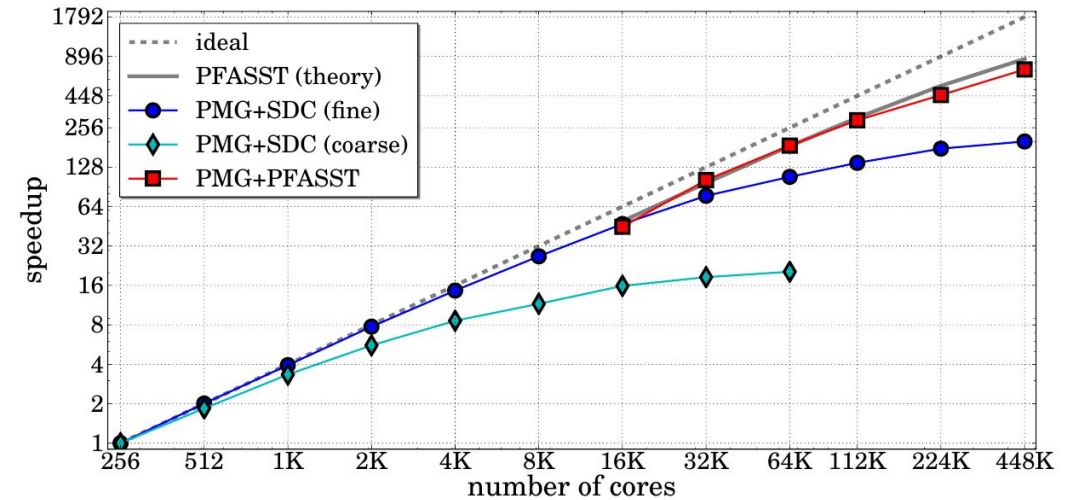


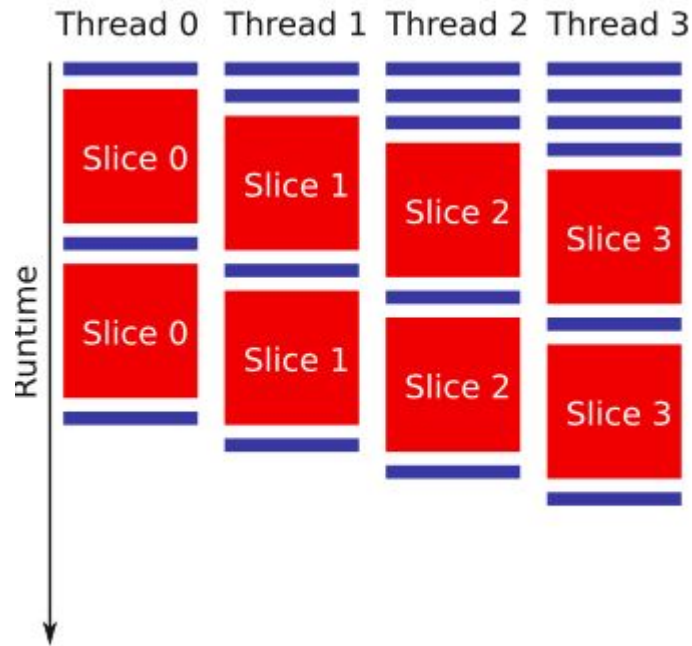
Figure courtesy T. Lunet

We can see this in practice! Sometimes.



* D. Ruprecht, R. Speck, M. Emmett, M. Bolten, and R. Krause, "Poster. Extreme-scale space-time parallelism," in Proceedings of the 2013 Conference on High Performance Computing Networking, Storage and Analysis Companion, ser. SC '13 Companion, Denver, Colorado, USA, 2013.

PinT: The challenge



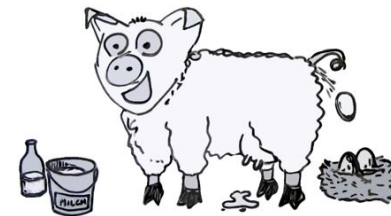
The **coarse propagator** is a serial bottleneck:

$$s(N_p) \leq \min \left(\frac{N_p}{N_{it}}, \frac{\text{runtime fine}}{\text{runtime coarse}} \right)$$

This is for *Parareal*, but similar bounds hold for other algorithms like *MGRIT* or *PFASSST*

Thus, coarse propagator needs to be both fast and reasonably accurate.

Proverbial “*Eier legende Wollmilchsau*” ... animal that lays eggs, gives milk and provides wool.



Using ML to build coarse propagators

$$f(V) = \frac{\partial V}{\partial t}(S, t) + \frac{1}{2}\sigma^2 S^2 \frac{\partial^2 V}{\partial S^2}(S, t) + rS \frac{\partial V}{\partial S}(S, t) - rV(S, t) = 0.$$

PDE Residual (used in loss function of PINN) risk-free interest rate
volatility

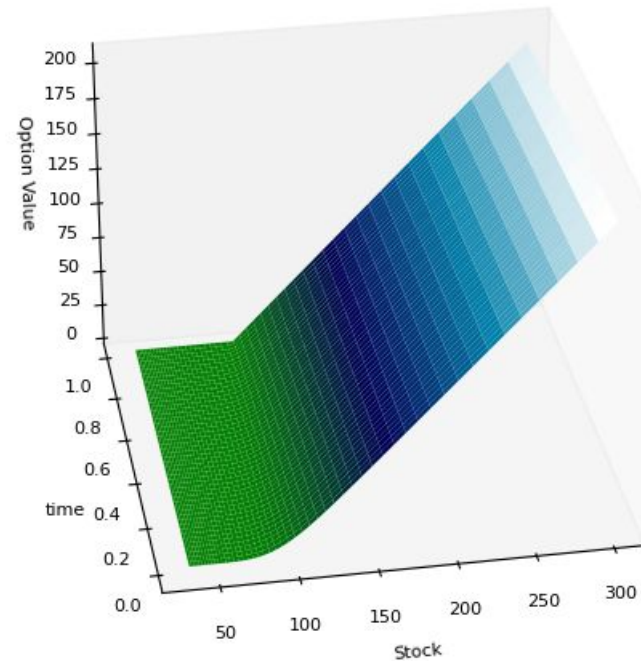
- Boundary conditions

$$V(t, S) \sim 0 \text{ as } S \rightarrow \infty, \text{ for fixed } t.$$

$$V(t, 0) = 0 \text{ for all } t.$$

- Expiration / end time condition

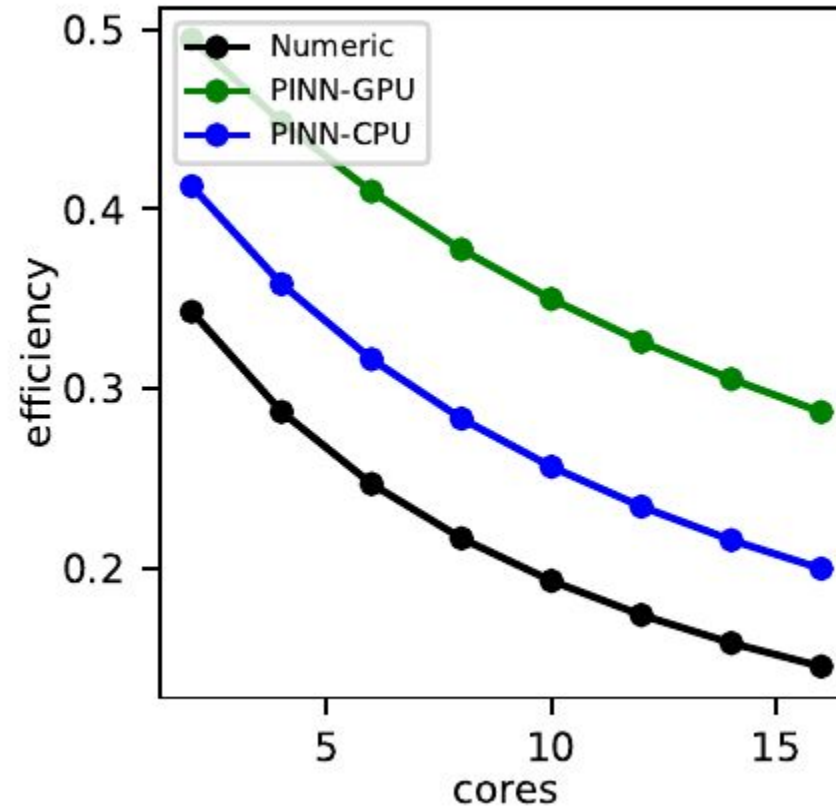
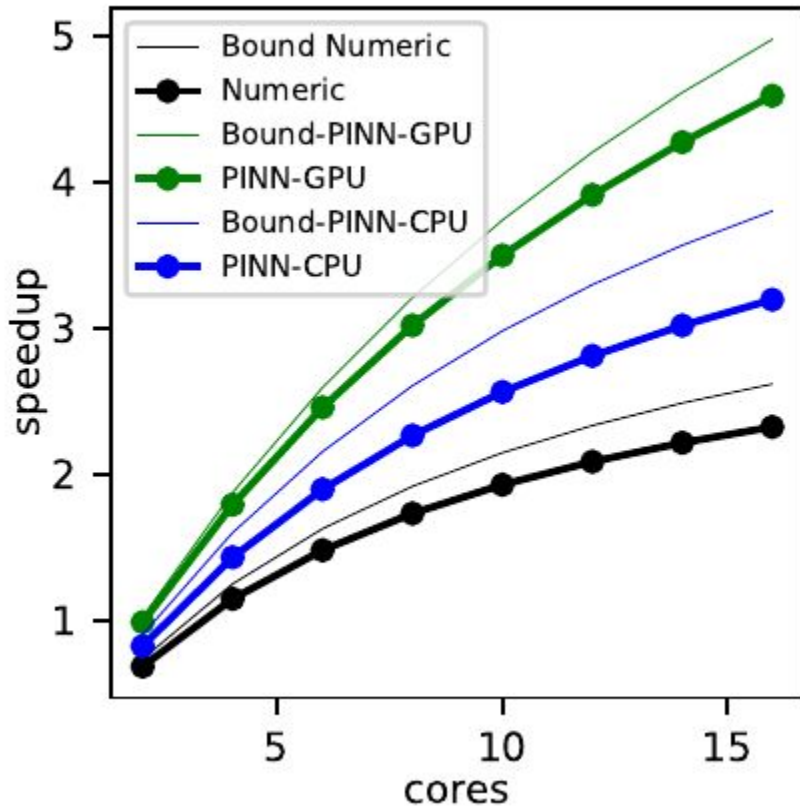
$$V(T, S) = \max(S - K, 0) \text{ for all } S$$



Work by: Abdul Qadir Ibrahim, M. Sc.

A. Q. Ibrahim, S. Götschel, and D. Ruprecht, "Parareal with a physics-informed neural network as coarse propagator," in **Euro-Par 2023: Parallel Processing**, Springer Nature Switzerland, 2023, pp. 649–663.

Using ML to build coarse propagators



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PINN coarse propagator on GPU yields more than twice the Parareal speedup than a mesh-based coarse propagator.

Performance modelling and comparison

$$\mathbf{u}_{n+1}^{k+1} = \mathbf{B}_0^0 \mathbf{u}_n^k + \mathbf{B}_1^0 \mathbf{u}_{n+1}^k + \mathbf{B}_0^1 \mathbf{u}_n^{k+1} + \dots$$

Algorithm	$\mathbf{B}_1^0 (\mathbf{u}_{n+1}^k)$	$\mathbf{B}_0^0 (\mathbf{u}_n^k)$	$\mathbf{B}_0^1 (\mathbf{u}_n^{k+1})$
damped Block Jacobi	$\mathbf{I} - \omega \mathbf{I}$	$\omega \phi^{-1} \chi$	–
ABJ	$\mathbf{I} - \tilde{\phi}^{-1} \phi$	$\tilde{\phi}^{-1} \chi$	–
ABGS	$\mathbf{I} - \tilde{\phi}^{-1} \phi$	–	$\tilde{\phi}^{-1} \chi$
PARAREAL	–	$(\phi^{-1} - \tilde{\phi}^{-1}) \chi$	$\tilde{\phi}^{-1} \chi$
TMG	$(1 - \omega)(\mathbf{I} - \mathbf{T}_C^F \phi_C^{-1} \mathbf{T}_F^C \phi)$	$\omega(\phi^{-1} - \mathbf{T}_C^F \phi_C^{-1} \mathbf{T}_F^C) \chi$	$\mathbf{T}_C^F \phi_C^{-1} \mathbf{T}_F^C \chi$
TMG _c	–	$(\phi^{-1} - \mathbf{T}_C^F \tilde{\phi}_C^{-1} \mathbf{T}_F^C) \chi$	$\mathbf{T}_C^F \tilde{\phi}_C^{-1} \mathbf{T}_F^C \chi$
TMG _f	$(\mathbf{I} - \mathbf{T}_C^F \phi_C^{-1} \mathbf{T}_F^C \phi)(\mathbf{I} - \tilde{\phi}^{-1} \phi)$	$(\tilde{\phi}^{-1} - \mathbf{T}_C^F \phi_C^{-1} \mathbf{T}_F^C \phi \tilde{\phi}^{-1}) \chi$	$\mathbf{T}_C^F \phi_C^{-1} \mathbf{T}_F^C \chi$
PFASST	$(\mathbf{I} - \mathbf{T}_C^F \tilde{\phi}_C^{-1} \mathbf{T}_F^C \phi)(\mathbf{I} - \tilde{\phi}^{-1} \phi)$	$(\tilde{\phi}^{-1} - \mathbf{T}_C^F \tilde{\phi}_C^{-1} \mathbf{T}_F^C \phi \tilde{\phi}^{-1}) \chi$	$\mathbf{T}_C^F \tilde{\phi}_C^{-1} \mathbf{T}_F^C \chi$

Can write and analyse different iterative PinT algorithms in a common framework (for linear problems)

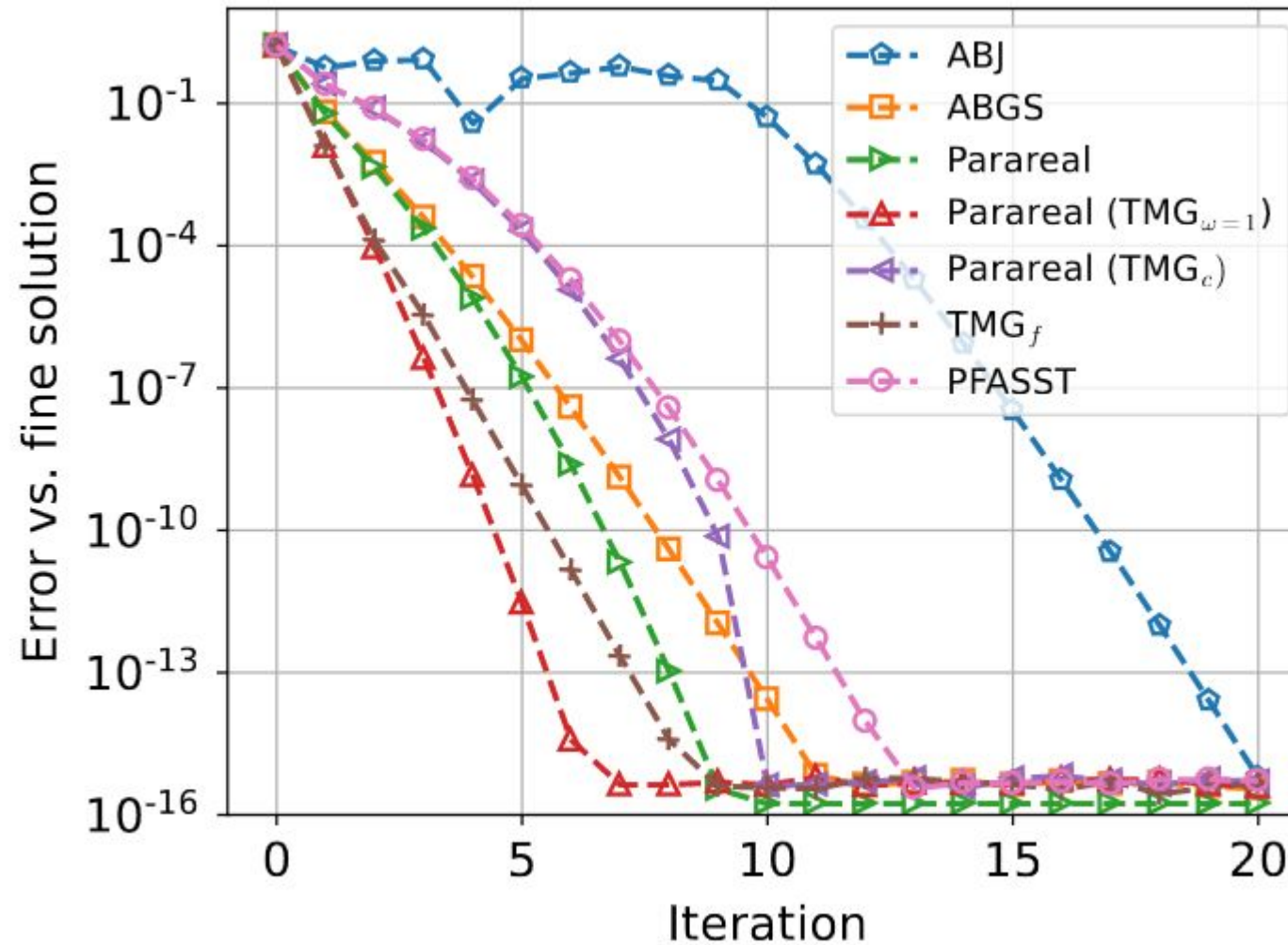


Work by: Dr
Thibaut Lunet.

M. J. Gander, T. Lunet, D. Ruprecht, and R. Speck, “A unified analysis framework for iterative parallel-in-time algorithms,” **SIAM Journal on Scientific Computing**, vol. 45, no. 5, 2275–A2303, 2023.

Performance modelling and comparison

Can predict
how
different
methods
will
converge



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Performance modelling and comparison

<https://time-x.eu/educational-website-for-first-performance-analysis-of-pint-algorithms/>

All required parameters set.

COMPUTE

Quick Documentation

Consider the Parareal block update formula :

$$u_{n+1}^k = F(u_n^k) + G(u_n^{k+1}) - G(u_n^k)$$

Define the number of blocks N (time-intervals), the number of iterations K (eventually different for each block), and the cost of F and G . Then choose a scheduler type :

- BLOCK-BY-BLOCK : uses N processors, where each processor is dedicated to one time block
- LOWEST-COST-FIRST : uses N processors, and compute first the tasks with lower cost
- OPTIMAL : eventually use more than N processors to minimize the overall computation time

Parareal Settings

N : 8
Number of iteration for each block (1 or N values)

K : 3

Computation time for F : 1

Computation time for G : 0.1

Scheduler Type : BLOCK-BY-BLOCK

Parareal Schedule

Processor rank: P7, P6, P5, P4, P3, P2, P1, P0

Time: 0, 1, 2, 3, 4

Performance Estimation ($N_p = 8$):

- Runtime : 4.1
- Speedup : 2.0
- Efficiency : 24.4%

FULL SCREEN



Work by: Dr
Thibaut Lunet.

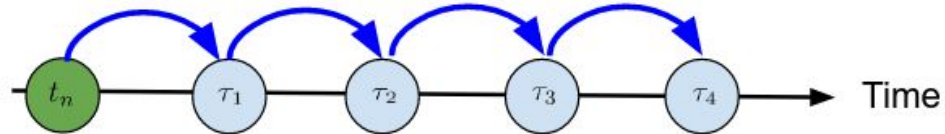
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Combined with a unified performance model by Bolten et al. 2023 from BU Wuppertal, we can predict performance of PinT methods on realistic HPC systems.

M. Bolten, S. Friedhoff, and J. Hahne, "Task graph-based performance analysis of parallel-in-time methods," *Parallel Computing*, vol. 118, p. 103050, 2023

Adaptivity and soft-fault resilience

Serial SDC

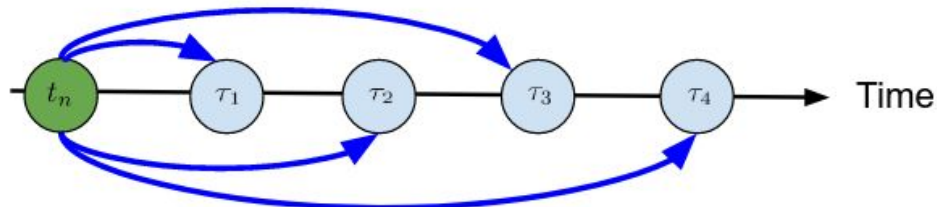


$$u_m^{k+1} = u_n + \Delta t \sum_{j=1}^M q_{m,j} f(u_j^k) + \Delta t \sum_{j=1}^m \Delta \tau_j [f(\tau_j, u_j^{k+1}) - f(\tau_j, u_j^k)]$$

► quadrature terms

► correction terms

Parallel SDC



$$u_m^{k+1} = u_n + \Delta t \sum_{j=1}^M q_{m,j} f(u_j^k) + \Delta t \alpha_m [f(\tau_m, u_m^{k+1}) - f(\tau_m, u_m^k)]$$

► quadrature terms

► correction terms

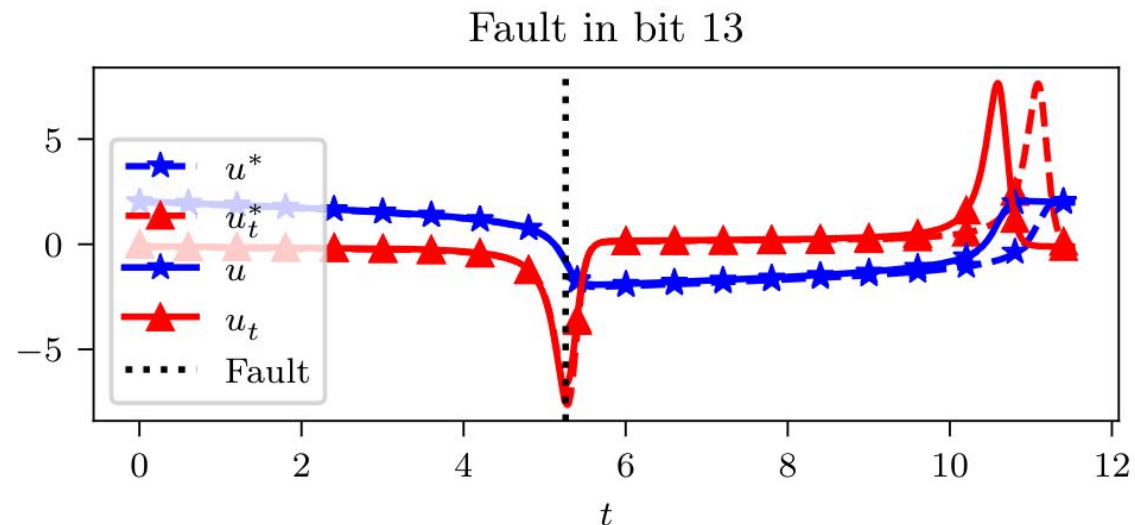
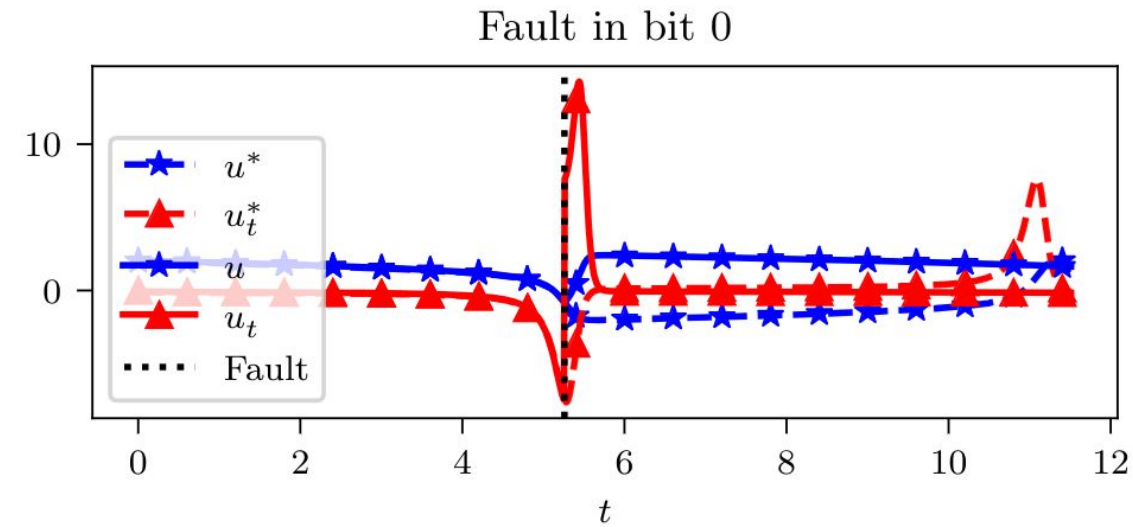


Work by: Thomas Baumann.

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Adaptivity and soft-fault resilience

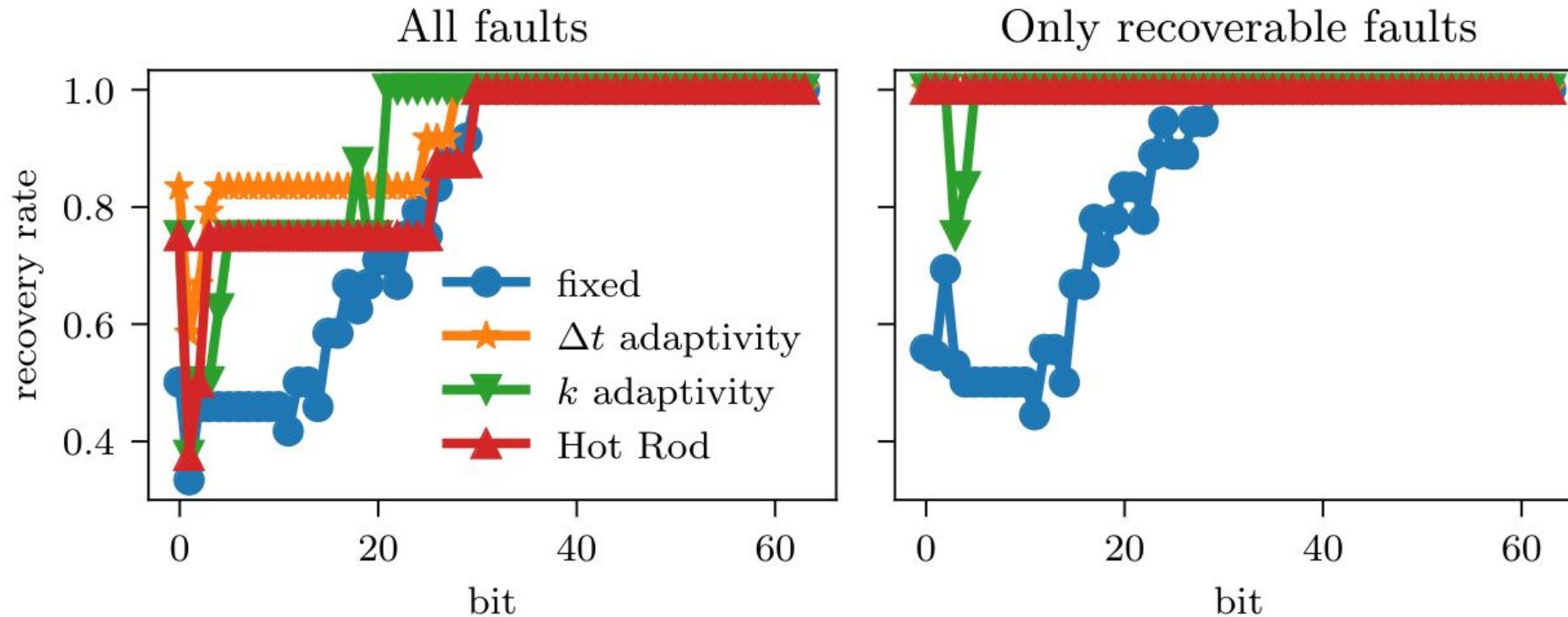
- Iterative nature of algorithm can protect against “bit flips”
- Flips show up in residual.
- Can mitigate, restart or continue to iterate and hope for the best!



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Adaptivity and soft-fault resilience



Suitable mitigation strategies can catch almost all faults in later bits and fix all recoverable faults.



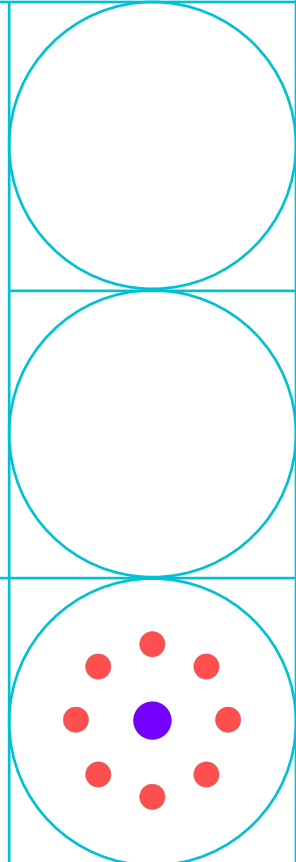
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Thanks.

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